

## Notes Toward Completion of the American Plan

Tony Roy – August 20, 2008

This document proposes a “completion” of the American Plan alternative to that in (3). It is striking that the difficulty for the four-valued approach occurs at precisely the same location as that addressed for the simplified semantics in (2). The current proposal is particularly natural against the background of solution there proposed.

### I. LANGUAGE / BASIC SEMANTIC NOTIONS

L4 The *vocabulary* consists of propositional parameters  $p_0, p_1 \dots$  with the operators  $\neg$ ,  $\wedge$ , and  $\rightarrow$ . Each propositional parameter is a *formula*; if  $A$  and  $B$  are formulas, so are  $\neg A$ ,  $(A \wedge B)$  and  $(A \rightarrow B)$ . Other operators abbreviate in the usual way. If  $A$  is a formula so formed, so is  $\overline{A}$ .

Let  $/A/$  and  $\backslash A \backslash$  represent either  $A$  or  $\overline{A}$  where what is represented is constant in a given context, but  $/A/$  and  $\backslash A \backslash$  are opposite. And similarly for other expressions with overlines as below.

I4 An *interpretation* for the basic logic DW is  $\langle W, N, \overline{N}, R, \overline{R}, h \rangle$  where  $W$  is a set of worlds;  $N, \overline{N} \subseteq W$  are normal worlds for truth and non-falsity respectively;  $R, \overline{R} \subseteq W^3$  are access relations for truth and non-falsity respectively; and  $h$  is a function which assigns 1 or 0 to each  $/p/$  at each  $w \in W$ . When  $h_w(/p/) = 1$  we say  $/p/$  *holds* at  $w$  and otherwise *fails*. As a constraint on interpretations we require also,

NC For any  $w \in /N/$ ,  $w/R/xy$  iff  $x = y$

Where  $x$  is empty or includes additional constraints as described below, a  $4x$  interpretation incorporates also any constraints in  $x$ .

H4 For complex expressions,

$(\neg)$   $h_w(/ \neg P /) = 1$  iff  $h_w(\backslash P \backslash) = 0$

$(\wedge)$   $h_w(/ P \wedge Q /) = 1$  iff  $h_w(/ P /) = 1$  and  $h_w(/ Q /) = 1$

$(\rightarrow)$   $h_w(/ P \rightarrow Q /) = 1$  iff there are no  $x, y \in W$  such that  $w/R/xy$  and either  $h_x(P) = 1$  but  $h_y(Q) = 0$ , or  $h_y(\overline{P}) = 1$  but  $h_x(\overline{Q}) = 0$

For a set  $\Gamma$  of formulas,  $h_w(\Gamma) = 1$  iff  $h_w(/ P /) = 1$  for each  $/ P / \in \Gamma$ ; then,

V4  $\Gamma \vDash_{4x} P$  iff there is no  $4x$  interpretation  $\langle W, N, \overline{N}, R, \overline{R}, h \rangle$  and  $w \in N$  such that  $h_w(\Gamma) = 1$  but  $h_w(P) = 0$ .

Other logics may be obtained by placing constraints on access. Numbering of conditions is kept mostly parallel to (1).

#### BASIC CONSTRAINTS

D3/4 If  $a/R/bx$  and  $xRcd$  then there is a  $y$  such that  $bRcy$  and  $a/R/yd$ , and a  $z$  such that  $bRzd$  and  $a/R/cz$ . And if  $a/R/xb$  and  $x\bar{R}cd$  then there is a  $y$  such that  $b\bar{R}cy$  and  $a/R/yd$ , and a  $z$  such that  $b\bar{R}zd$  and  $a/R/cz$ .

D5 If  $a/R/bc$  then there is a  $y$  such that  $a/R/by$  and  $yRbc$  and a  $z$  such that  $a/R/zc$  and  $z\bar{R}bc$ .

We have the effect of D20 ‘for free’ from the double condition from H4( $\rightarrow$ ).

PERMUTATION: As explained in (2) the standard simplified semantics requires a notion of *inclusion* in order to accommodate permutation. Insofar as we are working with a parallel version of the four-valued approach, something similar is required here. In this case, a family of inclusion notions is available and required.

$$\begin{aligned}
 a \leq b &\Rightarrow \begin{cases} \text{if } h_a(p) = 1 \text{ then } h_b(p) = 1 \text{ and if } h_b(\bar{p}) = 1 \text{ then } h_a(\bar{p}) = 1 \\ \text{if } bRxy \text{ then } aRxy \text{ if } a \notin N, \text{ otherwise if } b\bar{R}xy \text{ then } x \leq y \\ \text{if } a\bar{R}xy \text{ then } b\bar{R}xy \text{ if } b \notin \bar{N}, \text{ otherwise if } aRxy \text{ then } x \leq y \end{cases} \\
 a \leq^* b &\Rightarrow \begin{cases} \text{if } h_a(p) = 1 \text{ then } h_b(\bar{p}) = 1 \text{ and if } h_b(p) = 1 \text{ then } h_a(\bar{p}) = 1 \\ \text{if } b\bar{R}xy \text{ then } aRxy \text{ if } a \notin N, \text{ otherwise if } b\bar{R}xy \text{ then } x \leq y \\ \text{if } a\bar{R}xy \text{ then } bRxy \text{ if } b \notin N, \text{ otherwise if } a\bar{R}xy \text{ then } x \leq y \end{cases} \\
 a \leq^\sharp b &\Rightarrow \begin{cases} \text{if } h_a(\bar{p}) = 1 \text{ then } h_b(p) = 1 \text{ and if } h_b(\bar{p}) = 1 \text{ then } h_a(p) = 1 \\ \text{if } bRxy \text{ then } a\bar{R}xy \text{ if } a \notin \bar{N}, \text{ otherwise if } bRxy \text{ then } x \leq y \\ \text{if } aRxy \text{ then } b\bar{R}xy \text{ if } b \notin \bar{N}, \text{ otherwise if } aRxy \text{ then } x \leq y \end{cases}
 \end{aligned}$$

Given this we have,

D6 If  $aRbc$  then for some  $y \geq a$ ,  $bRyc$ , and for some  $z \geq^* a$ ,  $c\bar{R}bz$ . And if  $a\bar{R}bc$  then for some  $y \geq^\sharp a$ ,  $bRyc$ , and for some  $z \leq a$ ,  $c\bar{R}bz$

DW takes none of the extra constraints. TW has D3/4; RW D3/4 and D6; and R has D3/4, D5 and D6. Thus the American Plan is ‘completed’ at least through R.

## II. SOUNDNESS

BASIC PRINCIPLES: We show that the semantics is adequate for basic axioms and rules by direct arguments, a few of which are worked. Arguments are typically case heavy but quite parallel, and collapsed to some extent by the notation.

A1.  $A \rightarrow A$

- A2.  $A \rightarrow (A \vee B)$  and  $B \rightarrow (A \vee B)$
- A3.  $(A \wedge B) \rightarrow A$  and  $(A \wedge B) \rightarrow B$
- A4.  $[A \wedge (B \vee C)] \rightarrow [(A \wedge B) \vee C]$
- A5.  $[(A \rightarrow B) \wedge (A \rightarrow C)] \rightarrow [A \rightarrow (B \wedge C)]$
- A6.  $[(A \rightarrow C) \wedge (B \rightarrow C)] \rightarrow [(A \vee B) \rightarrow C]$
- A7.  $\neg\neg A \rightarrow A$
- R1.  $A, A \rightarrow B$  so  $B$
- R2.  $A, B$  so  $A \wedge B$
- R3.  $A \rightarrow B, C \rightarrow D$  so  $(B \rightarrow C) \rightarrow (A \rightarrow D)$
- C20.  $(A \rightarrow \neg B) \rightarrow (B \rightarrow \neg A)$

Suppose  $\not\vdash_4 (A \rightarrow \neg B) \rightarrow (B \rightarrow \neg A)$ ; then there is a  $\langle W, N, \overline{N}, R, \overline{R}, h \rangle$  and  $w \in N$  such that  $h_w(A \rightarrow \neg B) \rightarrow (B \rightarrow \neg A) = 0$ . Since  $w \in N$ , with NC there is some  $a \in W$  s.t.  $h_a(/A \rightarrow \neg B/) = 1$  and  $h_a(/B \rightarrow \neg A/) = 0$ . From the latter, there are  $x, y \in W$  s.t.  $a/R/xy$  and either  $h_x(B) = 1$  and  $h_y(\neg A) = 0$ , or  $h_y(\overline{B}) = 1$  and  $h_x(\overline{\neg A}) = 0$ . In the first case,  $h_y(\neg A) = 0$ ; so  $h_y(\overline{A}) = 1$ ; so with  $a/R/xy$  and  $h_a(/A \rightarrow \neg B/) = 1$ ,  $h_x(\overline{\neg B}) = 1$ ; so  $h_x(B) = 0$ ; which contradicts  $h_x(B) = 1$ . In the second case,  $h_x(\overline{\neg A}) = 0$ ; so  $h_x(A) = 1$ ; so with  $a/R/xy$  and  $h_a(/A \rightarrow \neg B/) = 1$ ,  $h_y(\neg B) = 1$ ; so  $h_y(\overline{B}) = 0$ ; which contradicts  $h_y(\overline{B}) = 1$ .

- C3.  $(A \rightarrow B) \rightarrow [(B \rightarrow C) \rightarrow (A \rightarrow C)]$  – given D3/4

Suppose  $\not\vdash_x (A \rightarrow B) \rightarrow [(B \rightarrow C) \rightarrow (A \rightarrow C)]$ ; then there is a  $\langle W, N, \overline{N}, R, \overline{R}, h \rangle$  and  $w \in N$  such that  $h_w((A \rightarrow B) \rightarrow [(B \rightarrow C) \rightarrow (A \rightarrow C)]) = 0$ . Since  $w \in N$ , with NC there is some  $a \in W$  s.t.  $h_a(/A \rightarrow B/) = 1$  and  $h_a(/(B \rightarrow C) \rightarrow (A \rightarrow C)/) = 0$ . From the latter, there are  $b, c \in W$  s.t.  $a/R/bc$  and either (i) or (ii). (i)  $h_b(B \rightarrow C) = 1$  and  $h_c(A \rightarrow C) = 0$ . From the latter, there are  $d, e \in W$  s.t.  $cRde$  and either  $h_d(A) = 1$  and  $h_e(C) = 0$ , or  $h_e(\overline{A}) = 1$  and  $h_d(\overline{C}) = 0$ . In the first case, from  $a/R/bc$  and  $cRde$  by D3/4 there is a  $z$  s.t.  $bRze$  and  $a/R/dz$ ; so with  $h_a(/A \rightarrow B/) = 1$  and  $h_d(A) = 1$ ,  $h_z(B) = 1$ ; so with  $h_b(B \rightarrow C) = 1$ ,  $h_e(C) = 1$  which contradicts  $h_e(C) = 0$ . In the second case, from  $a/R/bc$  and  $cRde$  by D3/4 there is a  $y$  s.t.  $bRdy$  and  $a/R/ye$ ; so with  $h_a(/A \rightarrow B/) = 1$  and  $h_e(\overline{A}) = 1$ ,  $h_y(\overline{B}) = 1$ ; so with  $h_b(B \rightarrow C) = 1$ ,  $h_d(\overline{C}) = 1$ ; which contradicts  $h_d(\overline{C}) = 0$ . (ii)  $h_c(\overline{B \rightarrow C}) = 1$  and  $h_b(A \rightarrow C) = 0$ . From the latter, there are  $d, e \in W$  s.t.  $b\overline{R}de$  and either  $h_d(A) = 1$  and  $h_e(C) = 0$ , or  $h_e(\overline{A}) = 1$  and  $h_d(\overline{C}) = 0$ . In the first case, from  $a/R/bc$  and  $b\overline{R}de$  by D3/4 there is a  $z$  s.t.  $c\overline{R}ze$  and  $a/R/dz$ ; so with  $h_a(/A \rightarrow B/) = 1$  and  $h_d(A) = 1$ ,  $h_z(B) = 1$ ; so with  $h_c(\overline{B \rightarrow C}) = 1$ ,  $h_e(C) = 1$

which contradicts  $h_e(C) = 0$ . In the second case, from  $a/R/bc$  and  $b\bar{R}de$  by D3/4 there is a  $y$  s.t.  $c\bar{R}dy$  and  $a/R/ye$ ; so with  $h_a(/A \rightarrow B/) = 1$  and  $h_e(\bar{A}) = 1$ ,  $h_y(\bar{B}) = 1$ ; so with  $h_c(\overline{B \rightarrow C}) = 1$ ,  $h_d(\bar{C}) = 1$ ; which contradicts  $h_d(\bar{C}) = 0$ .

C4.  $(A \rightarrow B) \rightarrow [(C \rightarrow A) \rightarrow (C \rightarrow B)]$  – given D3/4

C5.  $[A \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)$  – given D5

Suppose  $\not\vdash_x [A \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)$ ; then there is a  $\langle W, N, \bar{N}, R, \bar{R}, h \rangle$  and  $w \in N$  such that  $h_w([A \rightarrow (A \rightarrow B)] \rightarrow (A \rightarrow B)) = 0$ . Since  $w \in N$ , with NC there is some  $a \in W$  s.t.  $h_a(/A \rightarrow (A \rightarrow B)/) = 1$  and  $h_a(/A \rightarrow B/) = 0$ . From the latter, there are  $b, c \in W$  s.t.  $a/R/bc$  and either  $h_b(A) = 1$  and  $h_c(B) = 0$ , or  $h_c(\bar{A}) = 1$  and  $h_b(\bar{B}) = 0$ . Consider the first case. From  $a/R/bc$  with D5, there is a  $y$  such that  $a/R/by$  and  $yRbc$ ; so with  $h_a(/A \rightarrow (A \rightarrow B)/) = 1$  and  $h_b(A) = 1$ ,  $h_y(A \rightarrow B) = 1$ ; so  $h_c(B) = 1$ ; which contradicts  $h_c(B) = 0$ . Consider the second case. From  $a/R/bc$  with D5, there is a  $z$  such that  $a/R/zc$  and  $z\bar{R}bc$ ; so with  $h_a(/A \rightarrow (A \rightarrow B)/) = 1$  and  $h_c(\bar{A}) = 1$ ,  $h_z(\overline{A \rightarrow B}) = 1$ ; so  $h_b(\bar{B}) = 1$ ; which contradicts  $h_b(\bar{B}) = 0$ .

PERMUTATION: For this we need that the inclusion relations preserve generally conditions on atomics. As the arguments are similar, I work just the first.

L1 For any  $a, b \in W$ , if  $a \leq b$ , then (i) if  $h_a(P) = 1$ , then  $h_b(P) = 1$  and (ii) if  $h_b(\bar{P}) = 1$  then  $h_a(\bar{P}) = 1$ .

Suppose  $a \leq b$ ; then it is immediate that an atomic is such that (i) and (ii). So suppose that if  $a \leq b$  then  $A$  and  $B$  satisfy (i) and (ii).

Suppose  $P$  is  $\neg A$  and  $a \leq b$ . (i) Suppose  $h_a(\neg A) = 1$ ; then  $h_a(\bar{A}) = 0$ ; so by assp.  $h_b(\bar{A}) = 0$ ; so  $h_b(\neg A) = 1$ . (ii) Suppose  $h_b(\overline{\neg A}) = 1$ ; then  $h_b(A) = 0$ ; so by assp.  $h_a(A) = 0$ ; so  $h_a(\overline{\neg A}) = 1$ . And similarly for  $(\wedge)$ .

Suppose  $P$  is  $A \rightarrow B$  and  $a \leq b$ . (i) Suppose  $h_a(A \rightarrow B) = 1$ ; we want to show  $h_b(A \rightarrow B) = 1$ . Suppose otherwise, that  $h_b(A \rightarrow B) = 0$ ; then there are  $x, y \in W$  s.t.  $bRxy$  and either (1)  $h_x(A) = 1$  and  $h_y(B) = 0$ , or (2)  $h_y(\bar{A}) = 1$  and  $h_x(\bar{B}) = 0$ . We consider these in two cases: (a)  $a \notin N$ ; then from  $bRxy$  and  $a \leq b$  we have  $aRxy$ . (1)  $h_x(A) = 1$ ; so with  $h_a(A \rightarrow B) = 1$ , we have  $h_y(B) = 1$ , which contradicts  $h_y(B) = 0$ . (2)  $h_y(\bar{A}) = 1$ ; so with  $h_a(A \rightarrow B) = 1$ , we have  $h_x(\bar{B}) = 1$ , which contradicts  $h_x(\bar{B}) = 0$ . (b)  $a \in N$ ; then from  $bRxy$  and  $a \leq b$  we have  $x \leq y$ . (1)  $h_x(A) = 1$ ; so with  $a \in N$  and  $h_a(A \rightarrow B) = 1$ ,  $h_a(B) = 1$ ; so with  $x \leq y$  by assp.  $h_y(B) = 1$ , which contradicts  $h_y(B) = 0$ . (2)  $h_y(\bar{A}) = 1$ ; so with  $a \in N$  and  $h_a(A \rightarrow B) = 1$ ,  $h_y(\bar{B}) = 1$ ; so with  $x \leq y$  by assp.  $h_x(\bar{B}) = 1$ , which contradicts  $h_x(\bar{B}) = 0$ . And similarly for (ii).

- L2 For any  $a, b \in W$ , if  $a \leq^* b$ , then (i) if  $h_a(P) = 1$ , then  $h_b(\overline{P}) = 1$  and (ii) if  $h_b(P) = 1$  then  $h_a(\overline{P}) = 1$ .
- L3 For any  $a, b \in W$ , if  $a \leq^\# b$ , then (i) if  $h_a(\overline{P}) = 1$ , then  $h_b(P) = 1$  and (ii) if  $h_b(\overline{P}) = 1$  then  $h_a(P) = 1$ .

Now we are in a position to argue for C6 in the usual way.

C6.  $A \rightarrow [(A \rightarrow B) \rightarrow B]$  – given D6

Suppose  $\not\models_x A \rightarrow [(A \rightarrow B) \rightarrow B]$ ; then there is a  $\langle W, N, \overline{N}, R, \overline{R}, h \rangle$  and  $w \in N$  s.t.  $h_w(A \rightarrow [(A \rightarrow B) \rightarrow B]) = 0$ . Since  $w \in N$ , with NC there is some  $a \in W$  s.t. either (i) or (ii). (i)  $h_a(A) = 1$  and  $h_a((A \rightarrow B) \rightarrow B) = 0$ . From the latter, there are  $b, c \in W$  such that  $aRbc$  and either  $h_b(A \rightarrow B) = 1$  and  $h_c(B) = 0$ , or  $h_c(\overline{A \rightarrow B}) = 1$  and  $h_b(\overline{B}) = 0$ . Suppose the first of these; from  $aRbc$  by D6 there is a  $y \geq a$  s.t.  $bRyc$ ; from  $h_a(A) = 1$  and  $y \geq a$ , by L1,  $h_y(A) = 1$ ; so with  $bRyc$  and  $h_b(A \rightarrow B) = 1$ ,  $h_c(B) = 1$ , which contradicts  $h_c(B) = 0$ . Suppose the second, from  $aRbc$  by D6 there is a  $z \geq^* a$  s.t.  $c\overline{R}bz$ ; from  $h_a(A) = 1$  and  $z \geq^* a$ , by L2,  $h_z(\overline{A}) = 1$ ; so with  $c\overline{R}bz$  and  $h_c(\overline{A \rightarrow B}) = 1$ ,  $h_b(\overline{B}) = 1$ , which contradicts  $h_b(\overline{B}) = 0$ . (ii)  $h_a(\overline{A}) = 1$  and  $h_a((\overline{A \rightarrow B}) \rightarrow \overline{B}) = 0$ . From the latter, there are  $b, c \in W$  such that  $a\overline{R}bc$  and either  $h_b(A \rightarrow B) = 1$  and  $h_c(B) = 0$ , or  $h_c(\overline{A \rightarrow B}) = 1$  and  $h_b(\overline{B}) = 0$ . Suppose the first; from  $a\overline{R}bc$  by D6 there is a  $y \geq^\# a$  s.t.  $bRyc$ ; from  $h_a(\overline{A}) = 1$  and  $y \geq^\# a$ , by L3,  $h_y(A) = 1$ ; so with  $bRyc$  and  $h_b(A \rightarrow B) = 1$ ,  $h_c(B) = 1$ , which contradicts  $h_c(B) = 0$ . Suppose the second, from  $a\overline{R}bc$  by D6 there is a  $z \leq a$  s.t.  $c\overline{R}bz$ ; from  $h_a(\overline{A}) = 1$  and  $z \leq a$ , by L1,  $h_z(\overline{A}) = 1$ ; so with  $c\overline{R}bz$  and  $h_c(\overline{A \rightarrow B}) = 1$ ,  $h_b(\overline{B}) = 1$ , which contradicts  $h_b(\overline{B}) = 0$ .

### III. COMPLETENESS

Rather than show completeness directly, we set out to show that for any interpretation on the simplified semantics is a corresponding four-valued interpretation that preserves all the same truths. Completeness then follows directly from completeness on the simplified semantics.

BASIC SIMPLIFIED SEMANTICS: An *interpretation* is  $\langle W, g, R, \star, v \rangle$  where  $W$  is a set of worlds;  $g \in W$ ;  $R \subseteq W^3$ ;  $\star$  a function from  $W$  to  $W$ ; and  $v$  a function such that for any  $w \in W$  and  $p$ ,  $v_w(p) = 1$  or  $v_w(p) = 0$ . Let  $\leq$  be a reflexive, transitive relation on  $W$  such that if  $a \leq b$  then  $a \trianglelefteq b$  and  $b^\star \trianglelefteq a^\star$ , where,<sup>1</sup>

$$a \trianglelefteq b = \begin{cases} \text{if } v_a(p) = 1 \text{ then } v_b(p) = 1 \\ \text{if } bRxy \text{ and } a \neq g, \text{ then } aRxy \\ \text{if } bRxy \text{ and } a = g \text{ then } x \leq y \end{cases}$$

<sup>1</sup>Observe that this diverges from the definition in (1) and (2) where  $a \leq b$  requires  $b^\star \leq a^\star$ , which in turn requires  $a^{\star\star} \leq b^{\star\star}$ , etc. Given restrictions on  $\star$ , the accounts seem effectively the same, though the above above simplifies contact with other inclusion relations.

Then as constraints on interpretations, we require also, S: for any  $a \in W$ ,  $a = a^{**}$ ; NC: for any  $a, b \in W$ ,  $gRab$  iff  $a = b$ ; and D20: for any  $a, b, c \in W$ , if  $aRbc$  then  $aRc^*b^*$ . Where  $x$  is empty or includes additional constraints as described below, an  $Sx$  interpretation incorporates also any constraints in  $x$ .

$$(\neg) \ v_w(\neg A) = 1 \text{ iff } v_{w^*}(A) = 0$$

$$(\wedge) \ v_w(A \wedge B) = 1 \text{ iff } v_w(A) = 1 \text{ and } v_w(B) = 1$$

$$(\rightarrow) \ v_w(A \rightarrow B) = 1 \text{ iff there are no } x, y \in W \text{ such that } wRxy \text{ and } v_x(A) = 1 \text{ but } v_y(B) = 0$$

For a set  $\Gamma$  of formulas,  $v_w(\Gamma) = 1$  iff  $v_w(A) = 1$  for each  $A \in \Gamma$ ; then,

$$\text{VS}_x \ \Gamma \models_{Sx} A \text{ iff there is no } Sx \text{ interpretation } \langle W, g, R, \star, v \rangle \text{ such that } v_g(\Gamma) = 1 \text{ and } v_g(A) = 0.$$

OPTIONAL CONSTRAINTS: We require matched clauses for D5 and D6. In each case, these follow immediately the presence of D20 and S.

D3/4 If there is an  $x$  such that  $aRbx$  and  $xRcd$  then there is a  $y$  such that  $aRcy$  and  $bRyd$  and there is a  $z$  such that  $bRcz$  and  $aRzd$

D5 If  $aRbc$  then there is a  $y$  such that  $aRby$  and  $yRbc$  (and there is a  $z^*$  such that  $aRc^*z^*$  and  $z^*Rc^*b^*$ )

Suppose  $aRbc$ ; then by D20,  $aRc^*b^*$ ; so there is a  $y$  such that  $aRc^*y$  and  $yRc^*b^*$ ; set  $z = y^*$ ; then  $z^* = y^{**} = y$ ; so if  $aRbc$  there is a  $z^*$  such that  $aRc^*z^*$  and  $z^*Rc^*b^*$ .

D6 If  $aRbc$  then there is a  $y \geq a$  such that  $bRyc$ , (and there is a  $z^* \geq a$  such that  $c^*Rz^*b^*$ ).

Suppose  $aRbc$ ; then by D20,  $aRc^*b^*$ ; so there is a  $y \geq a$  such that  $c^*Ryb^*$ ; set  $z = y^*$ ; then  $z^* = y^{**} = y$ ; so if  $aRbc$ , there is a  $z^* \geq a$  such that  $c^*Rz^*b^*$ .

CONSTRUCTION AND RESULTS: For any  $\langle W, g, R, \star, v \rangle$  consider a corresponding  $\langle W, N, \bar{N}, R, \bar{R}, h \rangle$  such that there is a  $w \in W$  corresponding to each  $w \in W$ , where  $N = \{g\}$ ;  $\bar{N} = \{w \mid w^* = g\}$ ;  $R = \{\langle x, y, z \rangle \mid \langle x, y, z \rangle \in R\}$ ;  $\bar{R} = \{\langle x, y, z \rangle \mid \langle x^*, y, z \rangle \in R\}$ ;  $h_w(p) = v_w(p)$ ; and  $h_w(\bar{p}) = v_{w^*}(p)$ . And set  $a \leq b$  iff  $a \trianglelefteq b$  and  $b^* \trianglelefteq a^*$ ;  $a \leq^* b$  iff  $a \trianglelefteq b^*$  and  $b \trianglelefteq a^*$ ; and  $a \leq^\# b$  iff  $a^* \trianglelefteq b$  and  $b^* \trianglelefteq a$ .

L4 If  $\langle W, g, R, \star, v \rangle$  is an  $Sx$  interpretation then  $\langle W, N, \bar{N}, R, \bar{R}, h \rangle$  constructed as above is a  $4x$  interpretation such that if  $Dn \in x$  then  $Dn \in x$ .

Suppose  $w \in N$ ; then  $w = g$ . Say  $wRxy$ ; then  $gRxy$  and by construction,  $gRxy$ ; so by NC,  $x = y$ ; so  $x = y$ . Say  $x = y$ ; then  $x = y$ ; so by NC,  $gRxy$ ; so by construction,  $gRxy$ , which is to say  $wRxy$ . Suppose  $w \in \bar{N}$ ; then  $w^* = g$ . Say  $w\bar{R}xy$ ; then by

construction,  $w^*Rxy$ ; so  $gRxy$ ; so by NC,  $x = y$ ; so  $x = y$ . Say  $x = y$ ; then  $x = y$ ; so by NC,  $gRxy$ ; so  $w^*Rxy$ ; so by construction,  $w\bar{R}xy$ . So NC is satisfied.

Suppose  $a \leq^\# b$ ; then  $a^* \trianglelefteq b$  and  $b^* \trianglelefteq a$ . Say  $h_a(\bar{p}) = 1$ ; then  $v_{a^*}(p) = 1$  so by  $a^* \trianglelefteq b$ ,  $v_b(p) = 1$ ; so  $h_b(p) = 1$ . Suppose  $h_b(\bar{p}) = 1$ ; then  $v_{b^*}(p) = 1$ ; so by  $b^* \trianglelefteq a$ ,  $v_a(p) = 1$ ; so  $h_a(p) = 1$ . Suppose  $bRxy$  and  $a \notin \bar{N}$ ; then  $bRxy$  and  $a^* \neq g$ ; so by  $a^* \trianglelefteq b$ ,  $a^*Rxy$ ; so  $a\bar{R}xy$ . Suppose  $bRxy$  and  $a \in \bar{N}$ ; then  $bRxy$  and  $a^* = g$ ; so by  $a^* \trianglelefteq b$ ,  $x \leq y$ ; so  $x \leq y$ . Suppose  $aRxy$  and  $b \notin \bar{N}$ ; then  $aRxy$  and  $b^* \neq g$ ; so by  $b^* \trianglelefteq a$ ,  $b^*Rxy$ ; so  $b\bar{R}xy$ . Suppose  $aRxy$  and  $b \in \bar{N}$ ; then  $aRxy$  and  $b^* = g$ ; so by  $b^* \trianglelefteq a$ ,  $x \leq y$ ; so  $x \leq y$ . So  $a \leq^\# b$  has the right form; and similarly for  $a \leq b$  and  $a \leq^* b$ .

Suppose D3/4. Suppose  $aRbx$  and  $xRcd$ . Then by construction,  $aRbx$  and  $xRcd$ ; so by D3/4, there is a  $y$  such that  $aRcy$  and  $bRyd$  and there is a  $z$  such that  $bRcz$  and  $aRzd$ . So by construction, there is a  $y$  such that  $aRcy$  and  $bRyd$  and there is a  $z$  such that  $bRcz$  and  $aRzd$ . Suppose  $a\bar{R}bx$  and  $xRcd$ . Then by construction,  $a^*Rbx$  and  $xRcd$ ; so by D3/4, there is a  $y$  such that  $a^*Rcy$  and  $bRyd$  and there is a  $z$  such that  $bRcz$  and  $a^*Rzd$ . So by construction, there is a  $y$  such that  $a\bar{R}cy$  and  $bRyd$  and there is a  $z$  such that  $bRcz$  and  $a\bar{R}zd$ . These satisfy D3/4. And similarly in other cases.

Suppose D5. Suppose  $aRbc$ ; then by construction,  $aRbc$ ; so by D5, there is a  $y$  such that  $aRby$  and  $yRbc$ , and there is a  $z^*$  such that  $aRc^*z^*$  and  $z^*Rc^*b^*$ ; and with D20,  $aRzc$  and  $z^*Rbc$ ; so by construction, there is a  $y$  such that  $aRby$  and  $yRbc$ , and there is a  $z$  such that  $aRzc$  and  $z\bar{R}bc$ . Suppose  $a\bar{R}bc$ ; then by construction,  $a^*Rbc$ ; so by D5, there is a  $y$  such that  $a^*Rby$  and  $yRbc$ , and there is a  $z^*$  such that  $a^*Rc^*z^*$  and  $z^*Rc^*b^*$ ; and with D20,  $a^*Rzc$  and  $z^*Rbc$ ; so by construction, there is a  $y$  such that  $a\bar{R}by$  and  $yRbc$ , and there is a  $z$  such that  $a\bar{R}zc$  and  $z\bar{R}bc$ . In either case, D5 is satisfied.

Suppose D6. Suppose  $aRbc$ ; then  $aRbc$ . By D6 there is a  $y \geq a$  s.t.  $bRyc$ ; so  $bRyc$ ; and since  $y \geq a$ ,  $y \triangleright a$  and  $a^* \triangleright y^*$ ; so  $y \geq a$ . And by D6 again there is a  $z^* \geq a$  s.t.  $c^*Rz^*b^*$ ; and by D20,  $c^*Rbz$ ; so  $c\bar{R}bz$ ; and since  $z^* \geq a$ ,  $z^* \triangleright a$  and  $a^* \triangleright z^{**} = z$ ; so  $z \geq^* a$ . Suppose  $a\bar{R}bc$ ; then  $a^*Rbc$ . By D6 there is a  $y \geq a^*$  s.t.  $bRyc$ ; so  $bRyc$ ; and since  $y \geq a^*$ ,  $y \triangleright a^*$  and  $a = a^{**} \triangleright y^*$ ; so  $y \geq^\# a$ . By D6 again there is a  $z^* \geq a^*$  s.t.  $c^*Rz^*b^*$ ; and by D20,  $c^*Rbz$ ; so  $c\bar{R}bz$ ; and since  $z^* \geq a^*$ ,  $z^* \triangleright a^*$  and  $a = a^{**} \triangleright z^{**} = z$ ; so  $z \leq a$ . So D6 is satisfied.

L5 Where  $\langle W, N, R, \star, v \rangle$  and  $\langle W, N, \bar{N}, R, \bar{R}, v \rangle$  are as above, for any  $A$ , (i)  $h_w(A) = v_w(A)$  and (ii)  $h_w(\bar{A}) = v_{w^*}(A)$ .

Suppose  $\langle W, g, R, \star, v \rangle$  and  $\langle W, N, \bar{N}, R, \bar{R}, v \rangle$  are as above. By construction, atomics are such that (i) and (ii). So suppose,  $P$  and  $Q$  are such that (i) and (ii).

Suppose  $A$  is  $\neg P$ . (i)  $h_w(\neg P) = 1$  iff  $h_w(\overline{P}) = 0$ ; by assumption, iff  $v_{w^*}(P) = 0$ ; iff  $v_w(\neg P) = 1$ . (ii)  $h_w(\overline{\neg P}) = 1$  iff  $h_w(P) = 0$ ; by assumption, iff  $v_w(P) = 0$ ; iff  $v_{w^*}(\neg P) = 1$ . And similarly for  $\wedge$ .

Suppose  $A$  is  $P \rightarrow Q$ . (i) Suppose  $v_w(P \rightarrow Q) = 0$ ; then there are  $x, y \in W$  such that  $wRxy$  and  $v_x(P) = 1$  but  $v_y(Q) = 0$ ; so by assumption,  $h_x(P) = 1$  but  $h_y(Q) = 0$ . And since  $wRxy$ ,  $wRxy$ . So  $v_w(P \rightarrow Q) = 0$ . Suppose  $v_w(P \rightarrow Q) = 0$ ; then there are  $x, y \in W$  such that  $wRxy$  and either  $h_x(P) = 1$  but  $h_y(Q) = 0$ , or  $h_y(\overline{P}) = 1$  but  $h_x(\overline{Q}) = 0$ ; so by assumption, either  $v_x(P) = 1$  but  $v_y(Q) = 0$ , or  $v_{y^*}(P) = 1$  but  $v_{x^*}(Q) = 0$ . And from  $wRxy$ ,  $wRxy$ ; so in the first case,  $v_w(P \rightarrow Q) = 0$ . But for the second case, we have by D20, that  $wRy^*x^*$ , and so that  $v_w(P \rightarrow Q) = 0$ .

(ii) Suppose  $v_{w^*}(P \rightarrow Q) = 0$ ; then there are  $x, y \in W$  such that  $w^*Rxy$  and  $v_x(P) = 1$  but  $v_y(Q) = 0$ ; so by assumption,  $h_x(P) = 1$  but  $h_y(Q) = 0$ . And since  $w^*Rxy$ ,  $w\overline{R}xy$ . So  $v_w(P \rightarrow \overline{Q}) = 0$ . Suppose  $v_w(P \rightarrow \overline{Q}) = 0$ ; then there are  $x, y \in W$  such that  $w\overline{R}xy$  and either  $h_x(P) = 1$  but  $h_y(Q) = 0$ , or  $h_y(\overline{P}) = 1$  but  $h_x(\overline{Q}) = 0$ ; so by assumption, either  $v_x(P) = 1$  but  $v_y(Q) = 0$ , or  $v_{y^*}(P) = 1$  but  $v_{x^*}(Q) = 0$ . And since  $w\overline{R}xy$ ,  $w^*Rxy$ ; so in the first case,  $v_{w^*}(P \rightarrow Q) = 0$ . And for the second case, we have by D20, that  $w^*Ry^*x^*$ , and so that  $v_{w^*}(P \rightarrow Q) = 0$ .

COMPLETENESS: Suppose  $\Gamma \not\vdash_{\text{rx}} P$  for one of the relevant logics under consideration; then by the completeness of the simplified semantics,  $\Gamma \not\vdash_{\text{sx}} P$ ; so there is an  $Sx$  interpretation  $\langle W, g, R, \star, v \rangle$  s.t.  $v_g(\Gamma) = 1$  but  $v_g(P) = 0$ ; so by L4 and L5 there is a  $4x$  interpretation  $\langle W, N, \overline{N}, R, \overline{R}, v \rangle$  and  $w \in N$  such that  $h_w(\Gamma) = 1$  but  $h_w(P) = 0$ ; so  $\Gamma \not\vdash_{4x} P$ . So if  $\Gamma \vdash_{4x} P$  then  $\Gamma \vdash_{\text{rx}} P$ .

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